

# Planware II

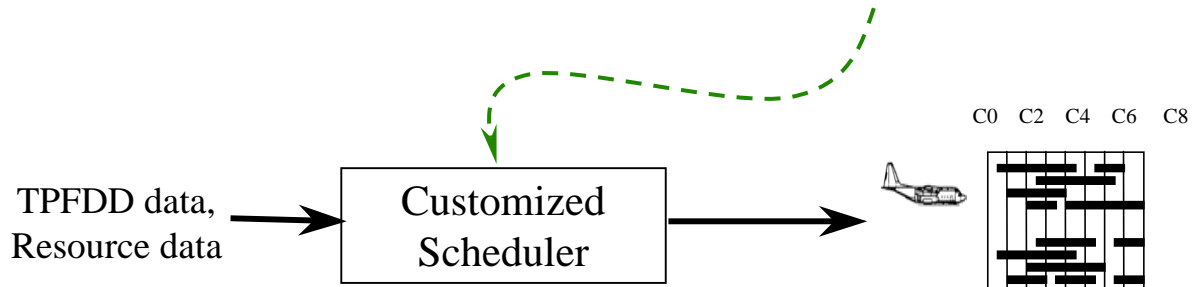
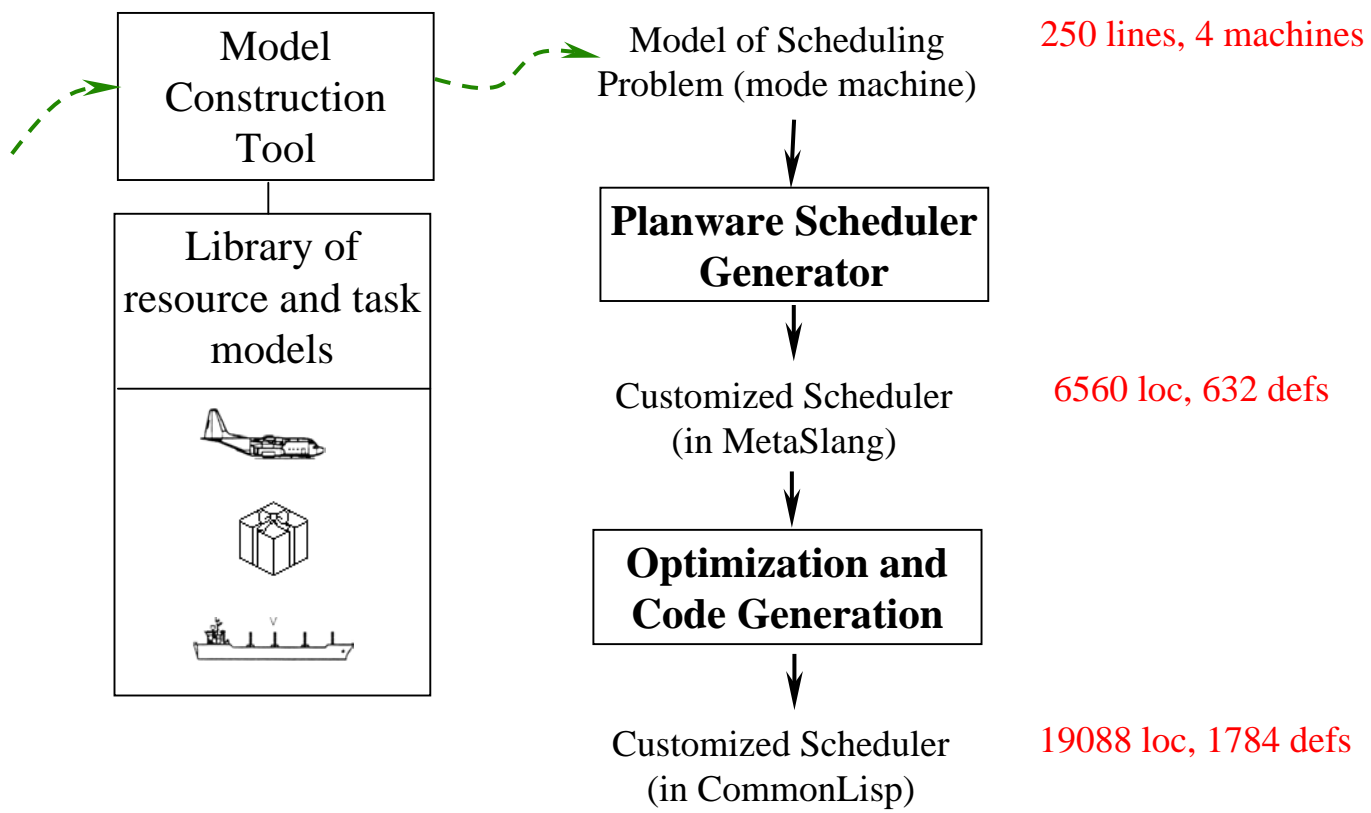
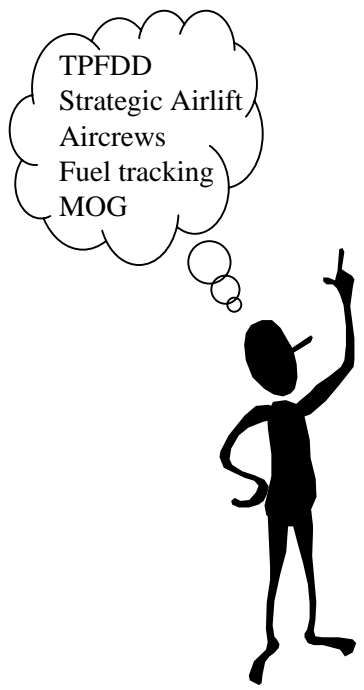
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# Planware: Synthesis of High Performance Schedulers



# How can we model resources?



What does a cargo aircraft do? e.g.

prep →fly →offload →prep →fly →fly →offload →fly →park →...

A *behavior* of a resource is a sequence of activities

Each *activity* is represented by a collection of attributes and their values

e.g. A “fly” activity of a cargo aircraft:

fly	
start-time :	2100 C0
finish-time :	0515 C1
duration :	8hr+15min
origin:	McGuire
destination:	Mildenhall
crew :	Jones
manifest:	{ ... }

A *resource* is characterized by its set of behaviors

e.g. cargo aircraft = { ...

prep →fly →offload →prep ...,

prep →fly →fly →offload →fly →park →... ,

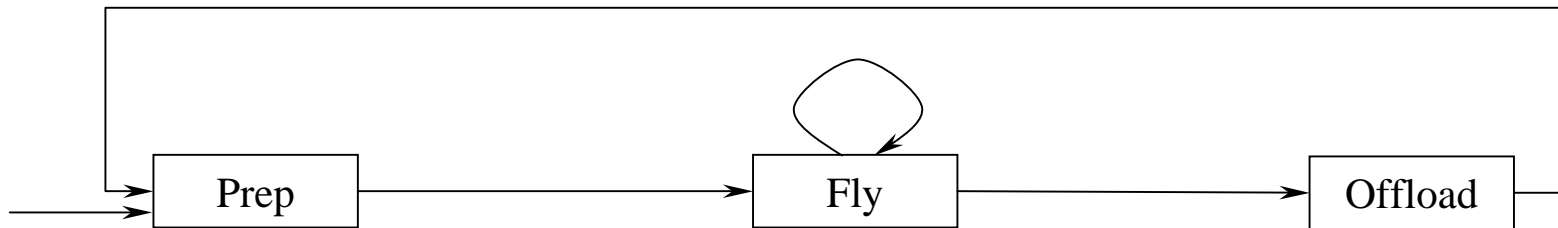
... }



# Key Idea:

## *model resources as state machines*

The (possibly infinite set of) behaviors of a resource can be expressed finitely by a state machine!



The *modes* of the state machine model the activities of the resource;  
The *transitions* model the restrictions on the sequence of activities

e.g. prep → fly → offload → prep → fly → fly → offload → prep → ...  
or prep → fly → fly → offload → prep → fly → fly → offload → prep → ...



# Modeling resources as state machines: *mode variables*

A signature presents the variables relevant to a resource:

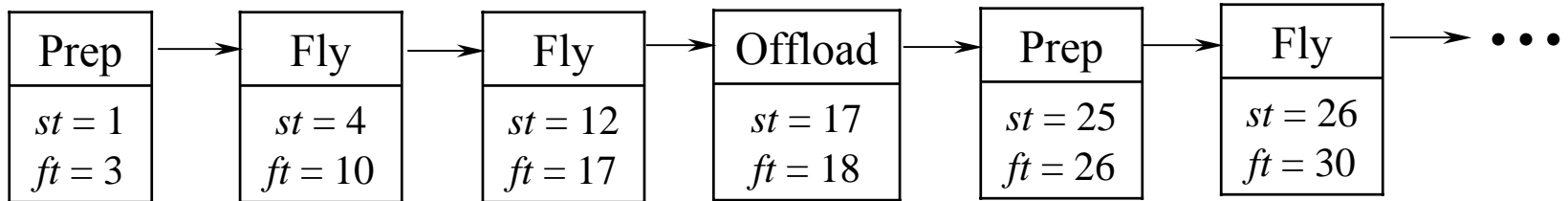
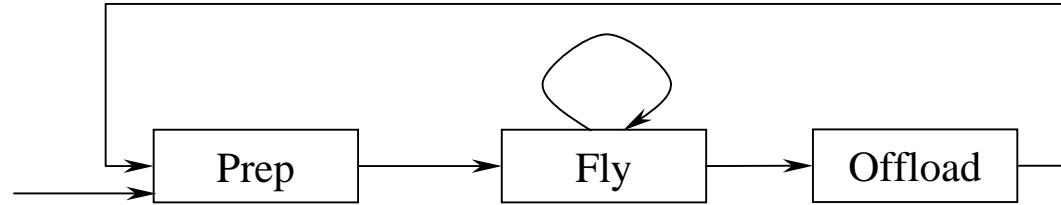
Signature for Cargo Aircraft	
var <i>st</i> : <i>Time</i>	– <i>start time</i>
var <i>ft</i> : <i>Time</i>	– <i>finish time</i>
var <i>dur</i> : <i>Duration</i>	– <i>duration</i>
var <i>orig</i> : <i>Port</i>	– <i>origin</i>
var <i>dest</i> : <i>Port</i>	– <i>destination</i>
var <i>mani</i> : <i>set(Task)</i>	– <i>manifest</i>
var <i>crew</i> : <i>Crew</i>	– <i>flight crew</i>
var <i>fuel-amt</i> : <i>Pounds</i>	– <i>fuel onload</i>
const <i>max-cap</i> : <i>Ston</i>	– <i>max acl</i>
var <i>loc</i> : <i>Location</i>	– <i>current position</i>



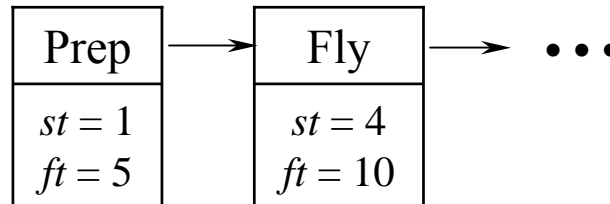
# Example Behaviors

Signature for Cargo Aircraft

```
var st : Time      - start time  
var ft : Time      - finish time  
...
```

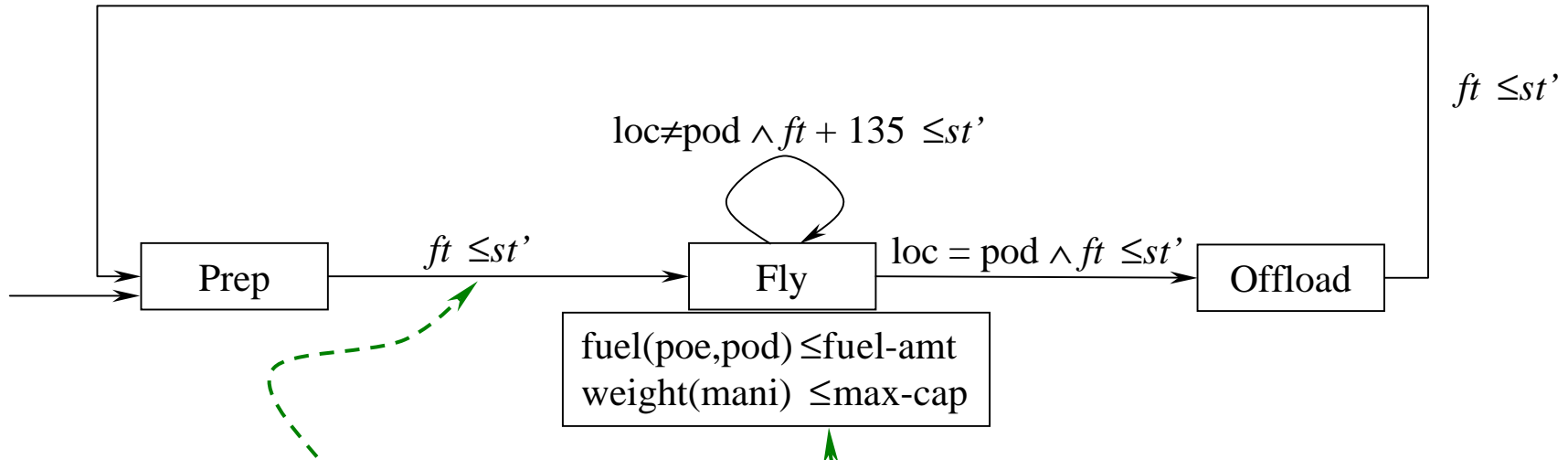


what about this?



# Modeling resources as state machines: *constraints*

Constraints on the modes and transitions constrain the evolution of the mode variables



The start time of the new activity is  
no earlier than the finish time of the previous activity

When flying, the resource must have adequate fuel  
and the weight of the manifest must not exceed capacity

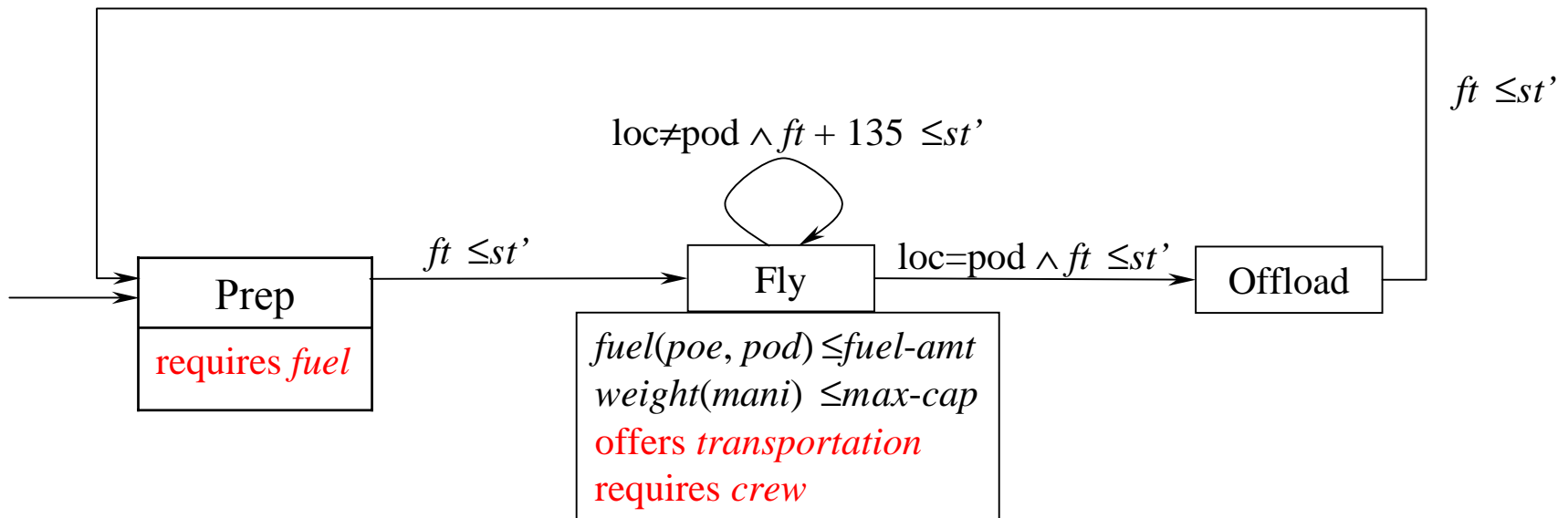
cf. in KIDS:  $\forall(ac:Aircraft, flt: Flight)$

$(ac \in sched \wedge flt \in ac.flights \Rightarrow weight(flt.mani) \leq max\text{-}cap)$



# Modeling resources as state machines: *services*

Each mode may *offer* services (to accomplish a certain kind of task)  
and each mode may *require* services of other resources





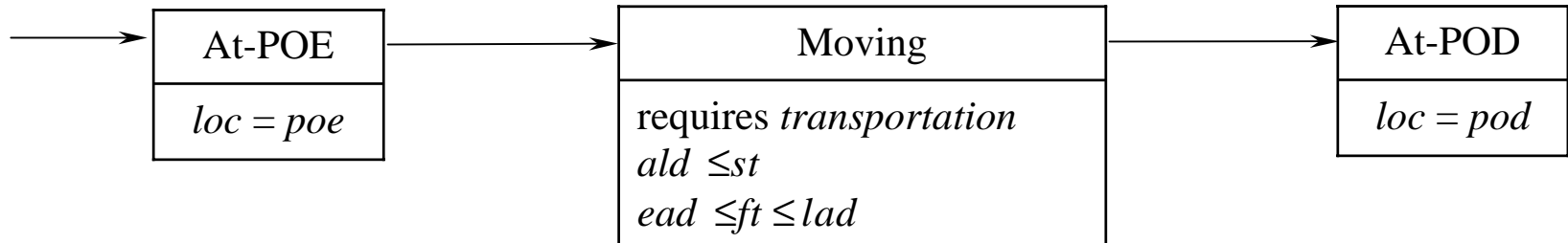
# Key Idea: *also model tasks as state machines!*

Effectively, a task is a resource that requires the services of other resources

e.g. a TPFDD-like  
movement requirement:

Signature for Move Req't.

```
var st, ft : Time  
var dur : Duration  
const poe, pod : Port  
const ald, ead, lad : Time  
const demand : Capacity  
var loc : Location
```

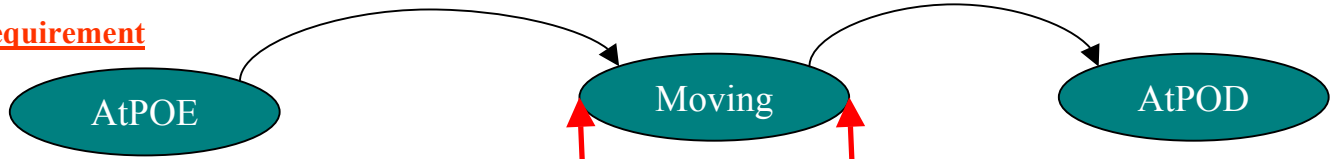


The scheduling process is motivated by the need to drive each task machine to a final/completed state, which entails scheduling the required resources, etc.

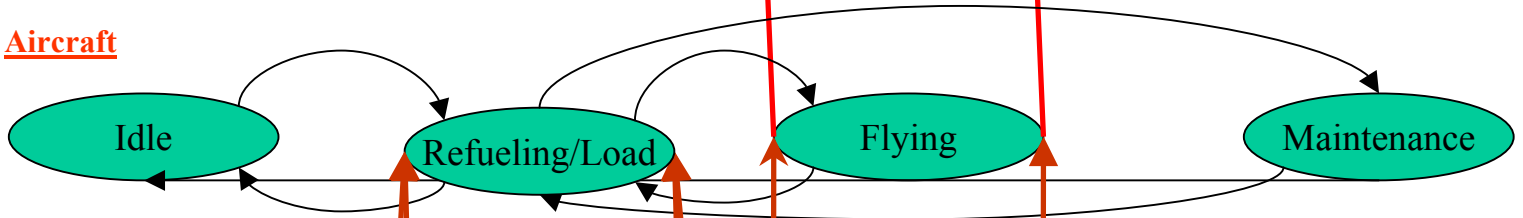


# Hierarchy of Required Services

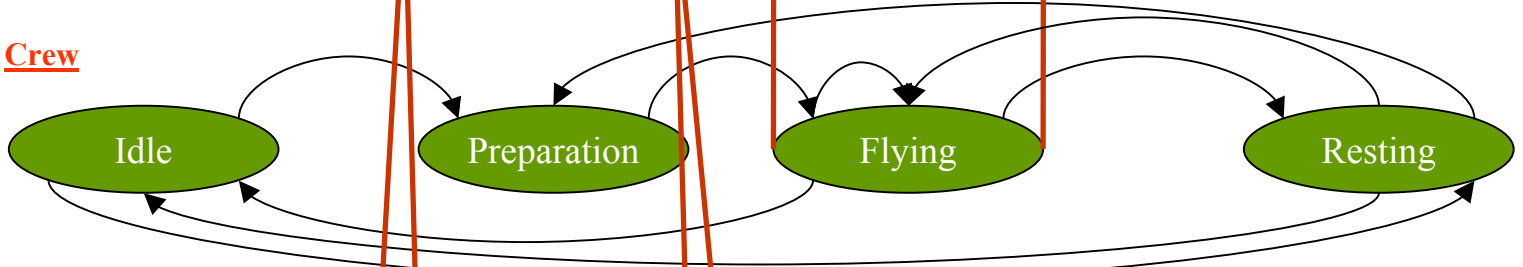
## MovementRequirement



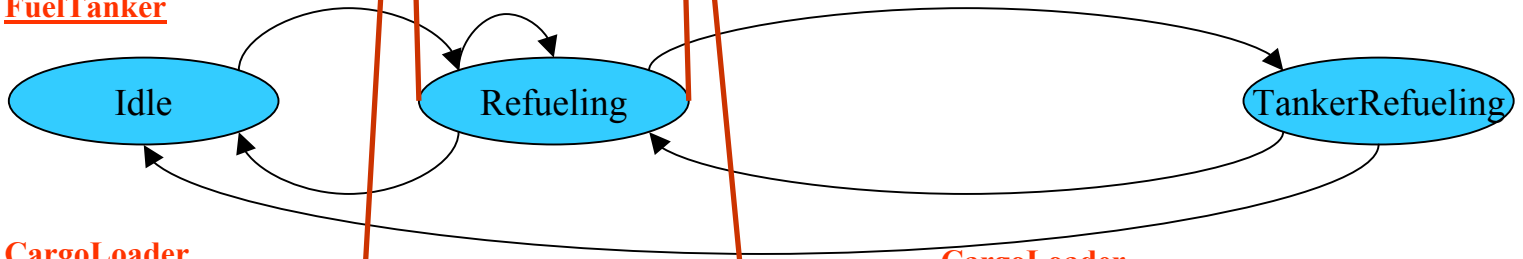
## Aircraft



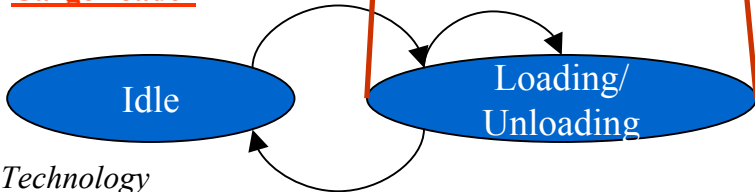
## Crew



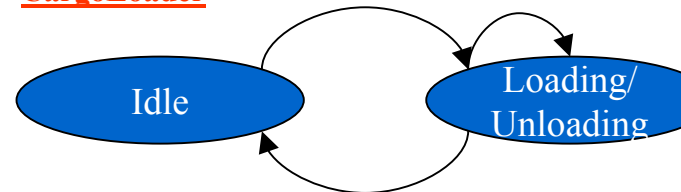
## FuelTanker



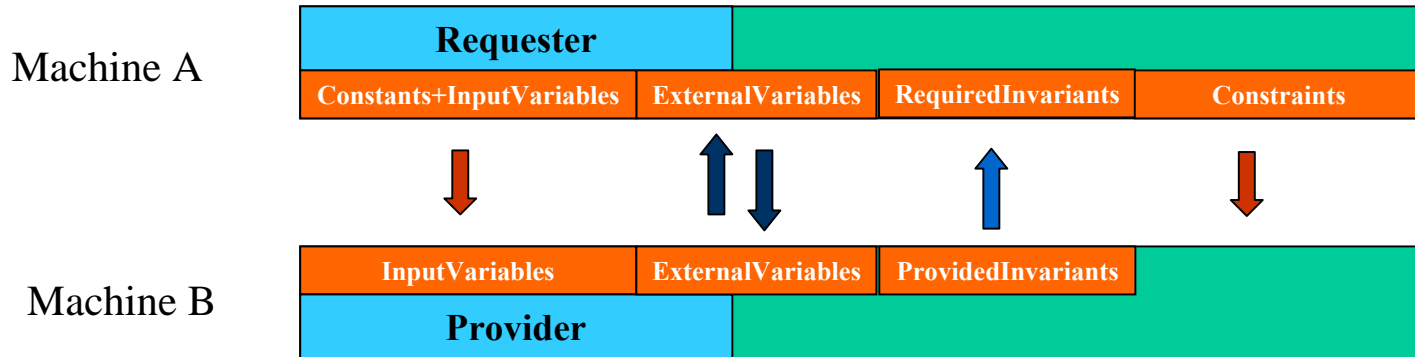
## CargoLoader



## CargoLoader



# Matching of Services



$$\begin{aligned}
 & \forall (\mathbf{constants(A)}, \mathbf{input-vars(A)}, \mathbf{constants(B)}) \\
 & \quad \exists (\mathbf{external-vars(A)}, \mathbf{internal-vars(A)}, \\
 & \quad \mathbf{input-vars(B)}, \mathbf{ext-vars(B)}, \mathbf{internal-vars(B)}) \\
 & (\mathbf{ProvidedConditions(A)}, \mathbf{ProvidedConditions(B)}) \\
 & \Rightarrow \\
 & (\mathbf{RequiredConditions(A)}, \mathbf{Constraints(A)}, \\
 & \mathbf{RequiredConditions(B)}, \mathbf{Constraints(B)})
 \end{aligned}$$



$$\forall(\text{constants}(A), \text{input-vars}(A), \text{constants}(B)) \exists(\text{ext-vars}(A), \text{internal-vars}(A), \text{input-vars}(B), \text{ext-vars}(B), \text{internal-vars}(B))$$

$$(\text{ProvidedConditions}(A), \text{ProvidedConditions}(B))$$

$$\Rightarrow$$

$$(\text{RequiredConditions}(A), \text{Constraints}(A), \text{RequiredConditions}(B), \text{Constraints}(B))$$

$\forall(\text{theMvr}, \text{POE}, \text{POD}, \text{ALD}, \text{LAD}, \text{EAD}, \text{theAc}, \text{maxFuel}, \text{maxCargoCap})$

$\exists(\text{ac}, \text{st}_{\text{mvr}}, \text{ft}_{\text{mvr}}, \text{dur}_{\text{mvr}}, \text{mvr}, \text{initialAcLocation}, \text{finalAcLocation}, \text{fuelLevel}, \text{st}_{\text{ac}}, \text{ft}_{\text{ac}}, \text{dur}_{\text{ac}},$   
 $\text{accFlyingHours}, \text{timeSinceMaint}, \text{currLocation}, \text{manifest}, \text{crew})$

$\text{POE} \neq \text{POD},$

$\text{CargoAtLocation}(\text{theMvr}, \text{POE}, \text{st}_{\text{mvr}}, \text{st}_{\text{mvr}}), \text{CargoAtLocation}(\text{theMvr}, \text{POD}, \text{ft}_{\text{mvr}}, \text{ft}_{\text{mvr}}),$

$\text{inTransit}(\text{mvr}, \text{theAc}, \text{initialAcLocation}, \text{finalAcLocation}, \text{st}_{\text{ac}}, \text{ft}_{\text{ac}}),$

$\text{inManifest}(\text{mvr}, \text{manifest}),$

$\Rightarrow$

$\text{inTransit}(\text{theMvr}, \text{ac}, \text{POE}, \text{POD}, \text{st}_{\text{mvr}}, \text{ft}_{\text{mvr}})$

$\text{st}_{\text{mvr}} \geq \text{ALD}, \text{ft}_{\text{mvr}} \geq \text{EAD}, \text{ft}_{\text{mvr}} \leq \text{LAD},$

$\text{initialAcLocation} \neq \text{finalAcLocation}, \text{atLocation}(\text{theAc}, \text{initialAcLocation}, \text{st}_{\text{ac}}, \text{st}_{\text{ac}}),$

$\text{hasFuel}(\text{theAc}, \text{fuelLevel}, \text{st}_{\text{ac}}, \text{st}_{\text{ac}}), \text{insideTransporter}(\text{mvr}, \text{theAc}, \text{st}_{\text{ac}}, \text{ft}_{\text{ac}}), \text{assignedCrew}(\text{theAc}, \text{crew}, \text{st}_{\text{ac}}, \text{ft}_{\text{ac}})$

$\text{timeDuration}(\text{ft}_{\text{ac}}, \text{st}_{\text{ac}}) \geq \text{transportDuration}(\text{theAc}, \text{initialAcLocation}, \text{finalAcLocation}),$

$\text{fuelLevel} > \text{requiredFuel}(\text{theAc}, \text{initialAcLocation}, \text{finalAcLocation})$

$\text{mvr} := \text{theMvr}, \text{ac} := \text{theAc}, \text{initialAcLocation} := \text{POE}, \text{finalAcLocation} := \text{POD}, \text{st}_{\text{mvr}} := \text{st}_{\text{ac}}, \text{ft}_{\text{mvr}} := \text{ft}_{\text{ac}},$

$\text{st}_{\text{ac}} \geq \text{ALD}, \text{ft}_{\text{ac}} \geq \text{EAD}, \text{ft}_{\text{ac}} \leq \text{LAD},$

$\text{atLocation}(\text{theAc}, \text{POE}, \text{st}_{\text{ac}}, \text{st}_{\text{ac}}),$

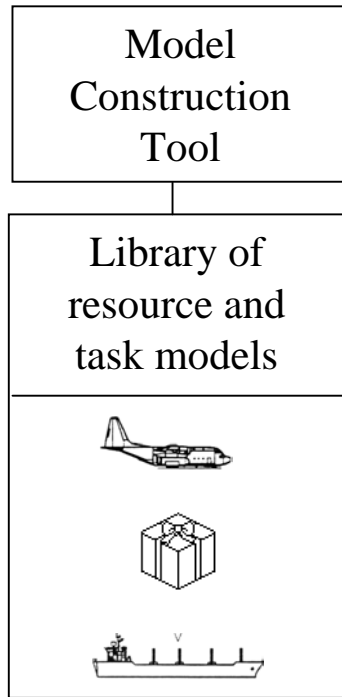
$\text{hasFuel}(\text{theAc}, \text{fuelLevel}, \text{st}_{\text{ac}}, \text{st}_{\text{ac}}), \text{insideTransporter}(\text{theMvr}, \text{theAc}, \text{st}_{\text{ac}}, \text{ft}_{\text{ac}}), \text{assignedCrew}(\text{theAc}, \text{crew}, \text{st}_{\text{ac}}, \text{ft}_{\text{ac}})$

$\text{timeDuration}(\text{ft}_{\text{ac}}, \text{st}_{\text{ac}}) \geq \text{transportDuration}(\text{theAc}, \text{initialLocation}, \text{finalLocation}),$

$\text{fuelLevel} > \text{requiredFuel}(\text{theAc}, \text{initialLocation}, \text{finalLocation}), \text{maxFuel} \geq \text{fuelLevel}$



# Code Generation Process



Model of Scheduling Problem  
(mode machine)

250 lines, 4 machines

**Planware Scheduler  
Generator**

Customized Scheduler  
(in MetaSlang)

6560 loc, 632 defs

**Optimization and  
Code Generation**

Customized Scheduler  
(in CommonLisp)

19088 loc, 1784 defs

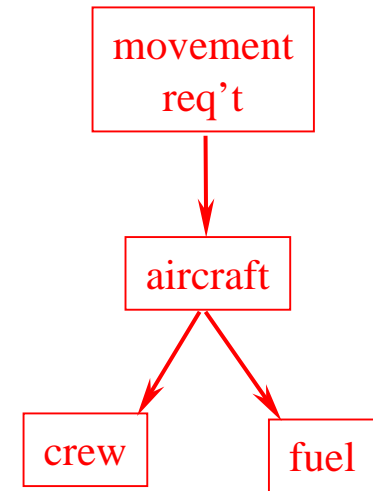
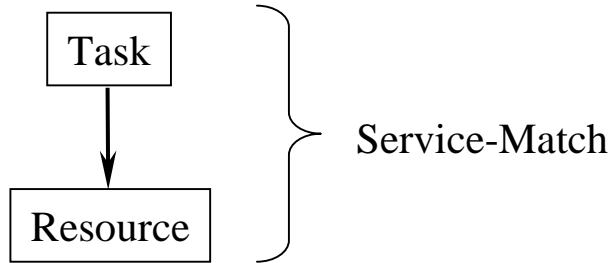


# Code Generation Statistics

Problem	Model LOC	MetaSlang LOC	CommonLisp LOC
Mvr, AC	161	3767	13,833
Mvr, AC, Crew	295	6167	20,225
Mvr, AC, Crew, Parking	444	8734	27,469
Mvr, AC missions, Crew, Port load/unload	536	12377	40,934



# Simplified Algorithm Schema



to schedule a task wrt a Service-Match:

get bids from all Resource instances;

consider the bids in sorted order

if the bid is feasible (via propagation)

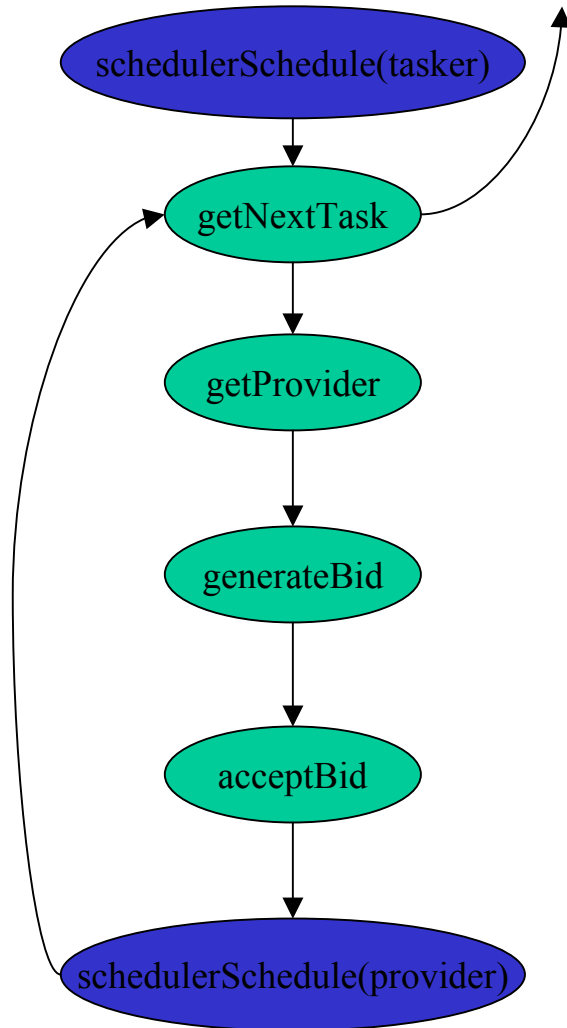
and the bid's resource generates a subtask

then schedule the subtask wrt its Service-Match

until one bid succeeds



# Program Schema I



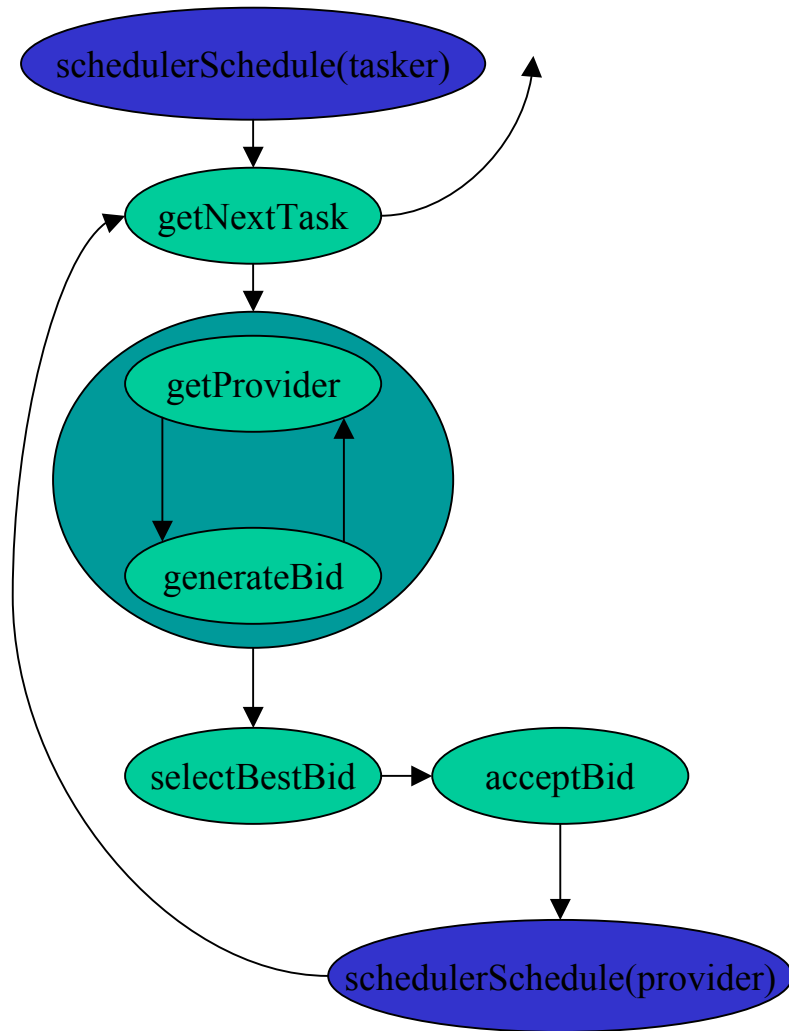
## Program Schema:

```
spec SchedulerSchedule(tasker: Tasker) =  
  while (nextTask := tasker.getNextTask())  
  do {  
    while (nextTask.unscheduled())  
    do {  
      provider := getNextProvider(nextTask)  
      if nextTask.satisfyRequirements(provider)  
      then {  
        bid := provider.generateBid(nextTask)  
  
        if feasible(bid)  
        then {  
          tasker.acceptBid(bid)  
          provider.AcceptBid(bid)  
          schedulerSchedule(provider)  
        }  
      }  
    }  
  }  
end-spec
```





# Program Schema II

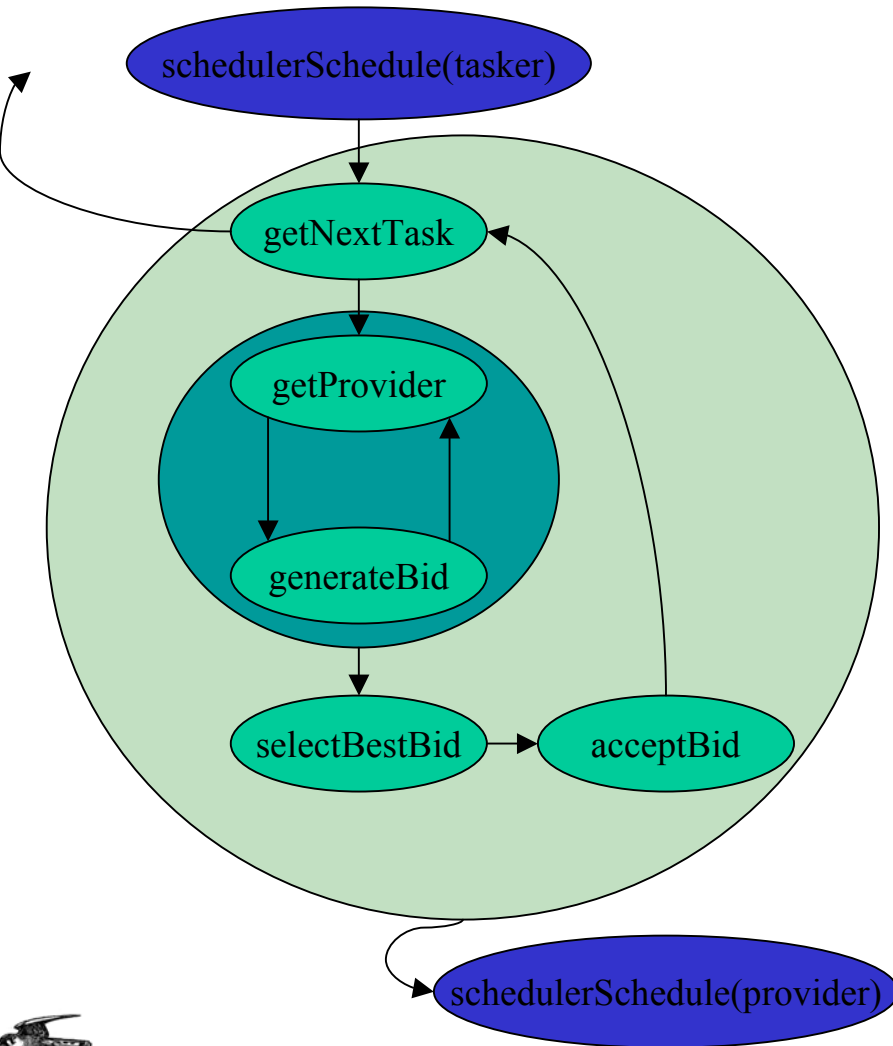


## Program Schema:

```
spec SchedulerSchedule(tasker: Tasker) =  
  while (nextTask := tasker.getNextTask())  
  do {  
    providers := findProvidersForTask(nextTask)  
    bids := []  
    for provider in providers  
    do {  
      if nextTask.satisfyRequirements(provider)  
      then {bids += provider.generateBid(nextTask)}  
    }  
    bestBid := getBestBid(bids)  
    resource := bestBid.getResource()  
    if feasible(bestBid)  
    then {  
      tasker.acceptBid(bestBid)  
      resource.AcceptBid(bestBid)  
      schedulerSchedule(resource)  
    }  
  }  
end-spec
```



# Program Schema III



## Program Schema:

```
spec SchedulerSchedule(tasker: Tasker) =  
  resources := []  
  while (nextTask := tasker.getNextTask())  
  do {  
    providers := findProvidersForTask(nextTask)  
    bids := []  
    for provider in providers  
    do {  
      if nextTask.satisfyRequirements(provider)  
      then { bids += provider.generateBid(nextTask)  
            }  
    }  
    bestBid := getBestBid(bids)  
    resource := bestBid.getResource()  
    if feasible(bestBid)  
    then {  
      tasker.acceptBid(bestBid)  
      resource.AcceptBid(bestBid)  
      resources += resource  
    }  
  }  
  for resource in resources  
  do { schedulerSchedule(resource) }  
end-spec
```



# Program Instance – Composing Program Schemas

```
program schedulerSchedule(tasker: VideoProcessor) =  
  while (nextTask := tasker.getNextTask())  
  do {  
    while (nextTask.unscheduled())  
    do {  
      provider := getNextProvider(nextTask)  
      if nextTask.satisfyRequirements(provider)  
      then {  
        bid := provider.generateBid(nextTask)  
        if feasible(bid)  
        then {  
          tasker.acceptBid(bid)  
          provider.AcceptBid(bid)  
          while (nextTask := provider.getNextTask())  
          do {  
            providers := findProvidersForTask(nextTask)  
            bids := []  
            for provider in providers  
            do {  
              if nextTask.satisfyRequirements(provider)  
              then {bids += provider.generateBid(nextTask) }  
            }  
            bestBid := getBestBid(bids)  
            resource := bestBid.getResource()  
            if feasible(bestBid)  
            then {  
              tasker.acceptBid(bestBid)  
              resource.AcceptBid(bestBid)  
              schedulerSchedule(resource)  
            }  
          }  
        }  
      }  
    }  
  }  
}
```



# Demonstration

## Steps in a doing a new project

- In Explorer: Create a new project folder
- In graphic window: build the machines (modes, transitions) and services
- In Explorer: designate initial and final modes
- In text window: for each machine, add constants, vars, constraints;  
add service parameters if necessary
- In text window: add initialization transitions  
fill out the other transitions
- In Explorer: generate service profile  
generate/compile/execute code



# Future Directions

- Modeling

- asynchronous servicing and subtasking via *multithreaded state machines!*
- improved integration with user-supplied components
- modeling objectives, preferences, priorities, and heuristics

- Code Generation

- improved optimizations, data structures, strategies
- algorithms for incremental rescheduling, distributed scheduling, online scheduling
- generate C, Java

- Planware environment

high quality total environment for modeling, scheduling, analysis, visualization, rescheduling



# Extras

