

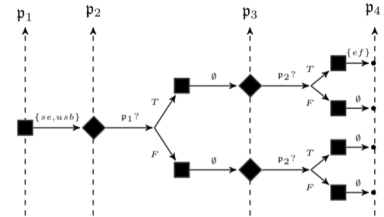
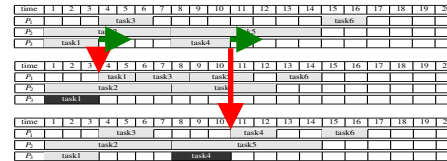
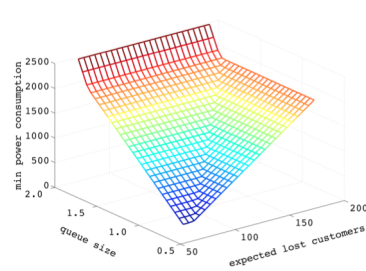
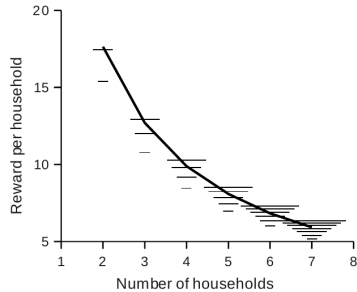


PRISM-games 3.0

Stochastic Game Verification

with

Concurrency, Equilibria and Time





PRISM-games 3.0

Stochastic Game Verification

with

Concurrency, Equilibria and Time

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University
of
Oxford

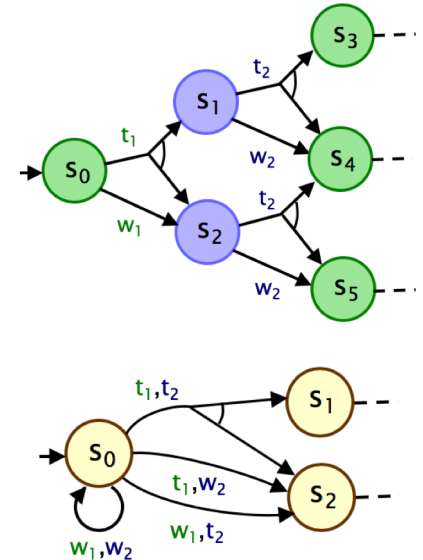
University
of
Glasgow

University
of
Birmingham

ERC
Advanced Grant
FUN2MODEL

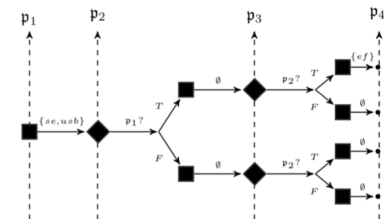
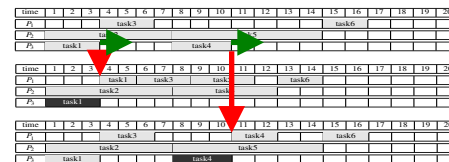
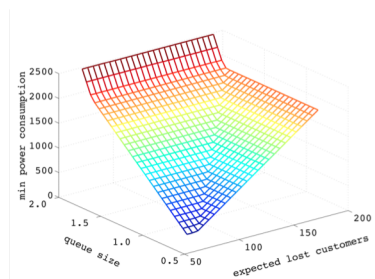
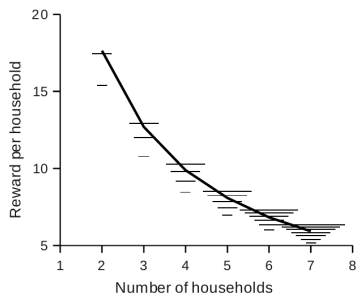
Stochastic Games

- Stochastic games:
 - **nondeterminism**: adversarial, controllers, ...
 - **probability**: randomisation, failures, noise, ...
 - **players**: competing/collaborating components
 - **strategies**: rational decisions made by players
 - **costs** (resources) & **rewards** (incentives)
- Applications:
 - computer security, network/communication protocols, algorithms for distributed consensus, energy management, autonomous robotics, auctions, ...
- Challenge:
 - how to design these systems correctly?
 - complex interactions between features



PRISM-games

- PRISM-games
 - **modelling** and analysis of **stochastic games**
 - automated **verification** or **synthesis** of strategies with **quantitative guarantees**
- Example specification in rPATL
 - $\langle\langle \text{robot}_1 \rangle\rangle P_{\geq 0.95} [F^{\leq 10} \text{goal}_1]$
 - “robot 1 has a strategy to ensure that, with probability at least 0.95, it reaches its goal in 10 steps, regardless of the strategies of other robots”



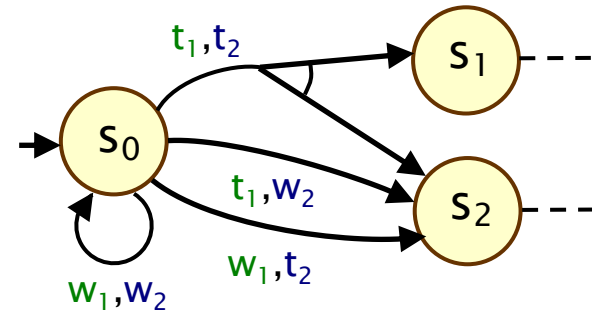
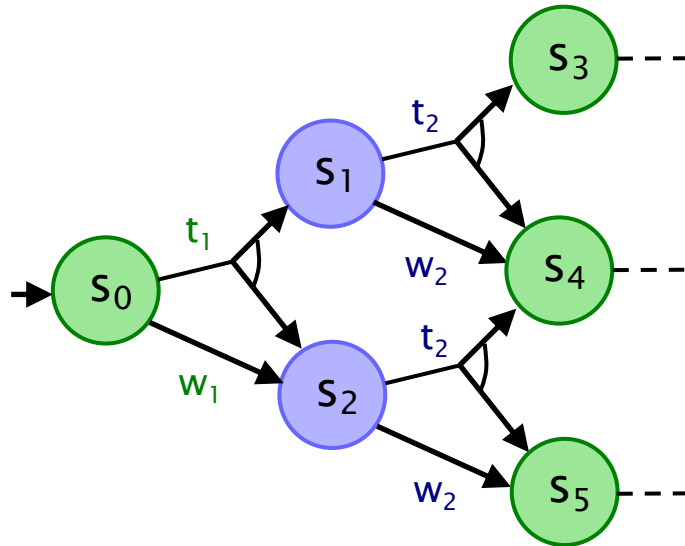
Concurrent stochastic games



Turn-based
stochastic games
(TSGs)



Concurrent
stochastic games
(CSGs)

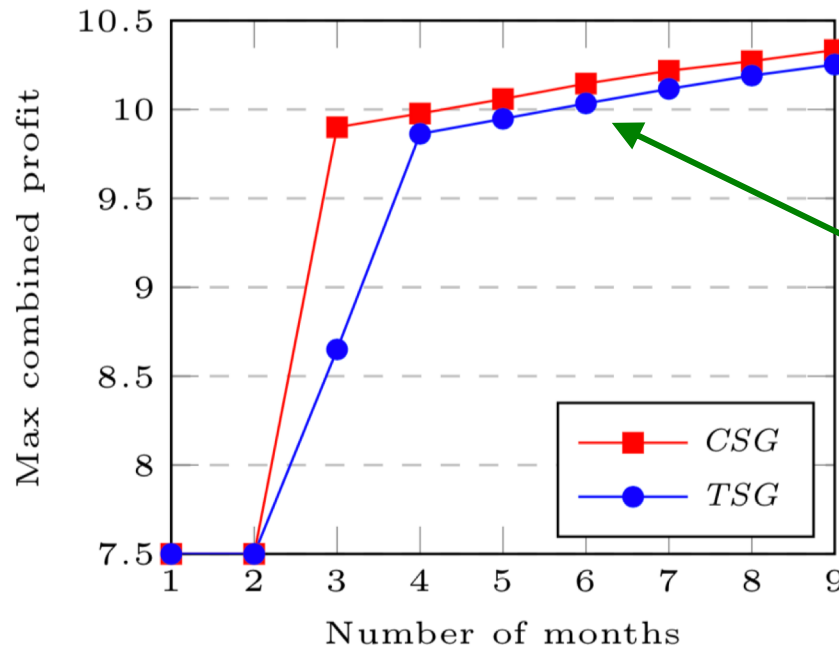


- Players make concurrent, independent decisions

Concurrent stochastic games



Example: Maximising expected investor profit in a futures market



Too pessimistic:
unrealistic strategy
for adversary

- Players make concurrent, independent decisions

CSGs in PRISM-games 3.0

```
csg
player p1 user1 endplayer
player p2 user2 endplayer
// Users (senders)
module user1
    s1 : [0..1] init 0; // has player 1 sent?
    e1 : [0..emax] init emax; // energy level of player 1
    [w1] true -> (s1'=0); // wait
    [t1] e1>0 -> (s1'=c' ? 0 : 1) & (e1'=e1-1); // transmit
endmodule
module user2 = user1 [ s1=s2, e1=e2, w1=w2, t1=t2 ] endmodule
// Channel: used to compute joint probability distribution for transmission failure
module channel
    c : bool init false; // is there a collision?
    [t1,w2] true -> q1 : (c'=false) + (1-q1) : (c'=true); // only user 1 transmits
    [w1,t2] true -> q1 : (c'=false) + (1-q1) : (c'=true); // only user 2 transmits
    [t1,t2] true -> q2 : (c'=false) + (1-q2) : (c'=true); // both users transmit
endmodule
```

CSGs in PRISM-games 3.0

csg

```
player p1 user1 endplayer  
player p2 user2 endplayer
```

```
// Users (senders)
```

```
module user1
```

```
  s1 : [0..1] init 0; // has player 1 sent?
```

```
  e1 : [0..emax] init emax; // energy level of player 1
```

```
  [w1] true -> (s1'=0); // wait
```

```
  [t1] e1>0 -> (s1'=c' ? 0 : 1) & (e1'=e1-1); // transmit
```

```
endmodule
```

```
module user2 = user1 [ s1=s2, e1=e2, w1=w2, t1=t2 ] endmodule
```

```
// Channel: used to compute joint probability distribution for transmission failure
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```
module channel
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  c : bool init false; // is there a collision?
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```

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  [w1,t2] true -> q1 : (c'=false) + (1-q1) : (c'=true); // only user 2 transmits
```

```
  [t1,t2] true -> q2 : (c'=false) + (1-q2) : (c'=true); // both users transmit
```

```
endmodule
```

Player = one or more modules

Variable updates
can refer to other
variables updated
simultaneously

Action lists
used to specify
synchronisation

Equilibria-based properties



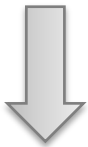
Zero-sum
properties



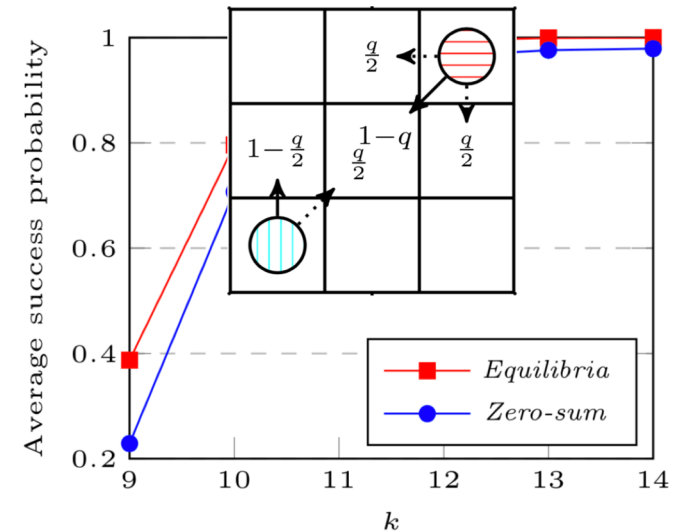
Equilibria-based
properties

- Example: multi-robot coordination

- $\langle\langle \text{robot}_1 \rangle\rangle_{\max=?} P [F^{\leq k} \text{goal}_1]$
- maximise probability of robot 1 reaching its goal, regardless of robot 2



- $\langle\langle \text{robot}_1 : \text{robot}_2 \rangle\rangle_{\max=?} (P [F^{\leq k} \text{goal}_1] + P [F^{\leq k} \text{goal}_2])$
- find strategies where robots 1 & 2 have no incentive to change actions and maximise joint goal probability



← Social-welfare optimal
Nash equilibrium

PRISM-games 3.0






- Probabilistic timed games (turn-based)



- 10 new/expanded case studies
 - multi-robot coordination, network trust models, Aloha, intrusion detection, public good games, ...



- More information at:
 - prismmodelchecker.org/games/
 - documentation, examples, case studies, papers
 - downloads:    + CAV'20 artefact VM



- Open source (GPLv2)
 - github.com/prismmodelchecker/prism-games