

$$\text{atig}(\text{Leaf}(n, b)) = \text{True}$$

$$\text{atig}(\text{Fork}(n, b) \ x \ y) =$$

$$\begin{aligned} & \text{if } b \text{ then } \text{btig } x \wedge \text{btig } y \\ & \text{else } \text{atig } x \wedge \text{atig } y \end{aligned}$$

$$\text{btig}(\text{Leaf}(n, b)) = \neg b$$

$$\text{btig}(\text{Fork}(n, b) \ x \ y) =$$

$$\neg b \wedge \text{atig } x \wedge \text{atig } y$$

$$\text{abtig } x = (\text{atig } x, \text{btig } x)$$

$$\text{abtig} = \text{foldTree base step}$$

$$\text{base}(n, b) = (\text{True}, \neg b)$$

$$\text{step}(n, b) \ (a_1, b_1) \ (a_2, b_2)$$

$$= (b \wedge b_1 \wedge b_2 \vee \neg b \wedge a_1 \wedge a_2, \neg b \wedge a_1 \wedge a_2)$$

$$\text{atig} = \text{fst} \cdot \text{abtig}$$

$$\text{mmp} : T \text{Int} \rightarrow \text{Int}$$

$$\text{mmp} = \max(\leq) \cdot \wedge (\text{value} \cdot \text{sat } p \cdot T \text{mark})$$

$$1) \quad \text{mark} : \text{Int} \xrightarrow{m} \text{Int} \times \text{Bool}$$

$$\text{mark } n = (n, \text{True}) \sqcup (n, \text{False})$$

$$2) \quad T : (X \xrightarrow{m} Y) \rightarrow (TX \xrightarrow{m} TY)$$

-defined later

$$3) \quad \text{sat} : (X \rightarrow \text{Bool}) \rightarrow (X \xrightarrow{m} X)$$

$$\text{sat } p = \text{ok} \cdot \langle \text{id}, p \rangle$$

$$\text{ok} (n, \text{True}) = n$$

4) $\text{value} : T(\text{Int} \times \text{Bool}) \rightarrow \text{Int}$

$\text{value} = \text{sum} \cdot T\text{val}$

$\text{val}(n, b) = \underline{\text{if } b \text{ then } n \text{ else } 0}$

$\text{sum} : T \text{Int} \rightarrow \text{Int}$

(defined later)

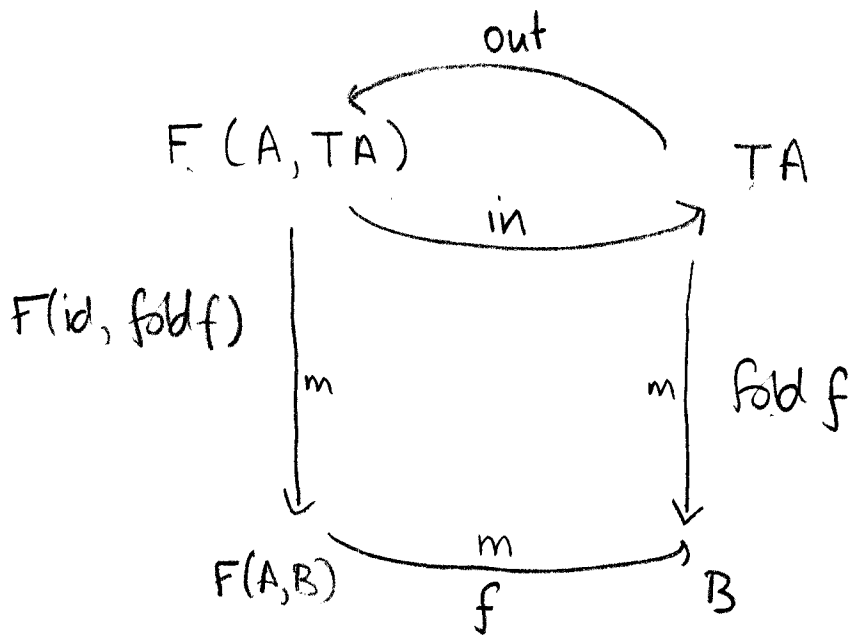
5) $\Lambda : (X \xrightarrow{m} Y) \rightarrow (X \rightarrow \text{Set } Y)$

$\Lambda f x = \{y \mid y \leftarrow f x\}$

6) $\text{max} : (X \rightarrow X \rightarrow \text{Bool}) \rightarrow (\text{Set } X \xrightarrow{m} X)$

$x \leftarrow \text{max} (\triangleleft) xs$ if

$x \in xs \wedge (\forall y \in xs : y \triangleleft x)$



$$id = fold\ in$$

$$Tf = fold\ (in \cdot F(f, id))$$

Functor fusion

$$fold\ f \cdot Tg = fold\ (f \cdot F(g, id))$$

p mutumorphism

$$p = prep \cdot fold\ \langle p_1, \dots, p_k \rangle$$

$$p_i : F(Int \times Bool, Bool^k) \rightarrow Bool$$

$$prep : Bool^k \rightarrow Bool$$

$$sum = fold\ plus$$

$$plus : F(Int, Int) \rightarrow Int$$

$$\begin{aligned}
& \text{value} \cdot \text{sat } p \\
= & \{ \text{above} \} \\
& \text{value} \cdot \text{ok} \cdot \langle \text{id}, p \rangle \\
= & \{ \text{claim} \} \\
& \text{ok} \cdot \langle \text{value}, p \rangle \\
= & \{ \text{defn} \} \\
& \text{ok} \cdot \langle \text{fold plus} \cdot \text{Tval}, p \rangle \\
= & \{ \text{functor fusion} \} \\
& \text{ok} \cdot \langle \text{fold} (\text{plus} \cdot F(\text{val}, \text{id})), \text{prop} \cdot \text{fold} \langle p_1 \dots p_k \rangle \rangle \\
= & \{ \text{op} = \text{ok} \cdot (\text{id} \times \text{prop}) \} \\
& \text{op} \cdot \langle \text{fold} (\text{plus} \cdot F(\text{val}, \text{id})), \text{fold} \langle p_1 \dots p_k \rangle \rangle \\
= & \{ \text{banana-split} \} \\
& \text{op} \cdot \text{fold } f
\end{aligned}$$

where $f = \langle \text{plus} \cdot F(\text{val}, \text{fst}), \langle p_1 \dots p_k \rangle \cdot F(\text{id}, \text{snd}) \rangle$

value · sat p · T mark

= {above}

op · fold f · T mark

= {functor fusion}

op · fold g

where $g = f \cdot F(\text{mark}, \text{id})$

Hence

$\max(\leq) \cdot \Lambda(\text{value} \cdot \text{sat } p \cdot \text{T mark})$

= {above}

$\max(\leq) \cdot \Lambda(\text{op} \cdot \text{fold } g)$

= {claim}

$\text{op} \cdot \max(\leq') \cdot \Lambda(\text{fold } g)$

where $(m, bs) \leq' (n, cs)$

$\hat{=} \text{prop } cs \wedge (\neg \text{prop } bs \vee m \leq n)$

↖ correction!

$(m, bs) \leq' (n, cs)$

$\equiv (\text{prop } bs, m) \leq_L (\text{prop } cs, n)$

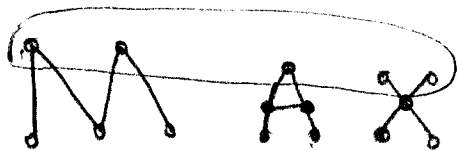
Thinning - a method for solving $\max(\leq) \cdot \wedge \text{fold } f$

$$\text{thin} : (X \rightarrow X \rightarrow \text{Bool}) \rightarrow (\text{Set } X \xrightarrow{m} \text{Set } X)$$

$$y_s \leftarrow \text{thin} (\triangleleft) x_s \text{ if}$$

$$y_s \subseteq x_s \wedge$$

$$(\forall x \in x_s : \exists y \in y_s : x \triangleleft y)$$



Theorem

$$\max(\leq) \cdot \wedge \text{fold } f \geq \max(\leq) \cdot \text{fold } g$$

$$\text{where } g : F(x, \text{Set } Y) \xrightarrow{m} \text{Set } Y$$

$$g = \text{thin} (\triangleleft) \cdot \wedge (f \cdot F(\text{id}, \text{choose}))$$

provided

$$(i) \quad x \triangleleft y \Rightarrow x \leq y$$

$$(ii) \quad f : F(x, Y) \xrightarrow{m} Y \text{ is monotonic}$$

under \triangleleft

$f: F(X, Y) \xrightarrow{m} Y$ monotonic

under (\leq) : $Y \rightarrow Y \rightarrow \text{Bool}$

if $x \leq_F y \wedge u \leftarrow fx$

then $(\exists v: v \leftarrow fy: u \leq v)$

$(\leq_F): F(X, Y) \rightarrow F(X, Y) \rightarrow \text{Bool}$

For our problem:

$f = \langle \text{plus} \cdot F(\text{val} \cdot \text{mark}, \text{fst}), \langle P_1 \dots P_k \rangle \cdot F(\text{mark}, \text{snd}) \rangle$

and it is monotonic under \leq where

$(m, bs) \leq (n, cs)$

$\hat{=} (m \leq n \wedge bs = cs)$

$(m, bs) \leq (n, cs)$

$\Rightarrow (\text{prop } bs, m) \leq_L (\text{prop } cs, n)$

In thinning $\text{mba} \leq$, we need keep no more than 2^k partial solutions.

Where is the program?

$mmp = \text{maxlist } r \cdot \text{fold} (\text{thinlist } q \cdot \text{extend})$

$\text{extend} = \text{concat} \cdot \text{map } \text{lambdaf} \cdot \text{cplist}$

$\text{cplist} : F(A, \text{List } B) \rightarrow \text{List } (F(A, B))$

$\text{cplist} \in \text{listify} \cdot \wedge (F(\text{id}, \text{choose}))$

$\text{lambdaf} : F(\text{Int}, \text{Int} \times \text{Bool}^k) \rightarrow \text{List } (\text{Int} \times \text{Bool}^k)$

$\text{lambdaf} \in \text{listify} \cdot \wedge f$

plug in F , plus, $P_1 \dots P_k$,

and play!

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