

Exercise Sheet 3 for Categories, Proofs and Games

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1. The notion of *universal arrow* is dual to that of couniversal arrow. We state it explicitly. Let $G : \mathcal{D} \longrightarrow \mathcal{C}$ be a functor, and A an object of \mathcal{C} . A *universal arrow from A to G* is an object D of \mathcal{D} and a morphism $f : A \longrightarrow GD$ such that, for every object D' of \mathcal{D} , and morphism $g : A \longrightarrow GD'$, there exists a unique morphism $h : D \longrightarrow D'$ such that $g = Gh \circ f$.

- Show carefully that a coproduct of objects A, B of \mathcal{C} is a universal arrow to the diagonal functor

$$\Delta : \mathcal{C} \longrightarrow \mathcal{C} \times \mathcal{C}.$$

- Let $U : \mathbf{Mon} \longrightarrow \mathbf{Set}$ be the ‘forgetful’ functor which sends a monoid to its set of elements. Show that for each set X , there is a universal arrow $X \longrightarrow UX^*$, where X^* is the monoid of finite sequences of elements of X , with concatenation as the binary operation.
 - (If you have not encountered rings, you should skip this part). Let \mathbf{Ring} be the category with commutative rings with unit as objects, and ring homomorphisms as morphisms. Let \mathbf{Ring}_* be the category where objects are pairs (R, a) where R is a ring, and $a \in R$; the morphisms from (R, a) to (S, b) are the ring homomorphisms $h : R \longrightarrow S$ such that $h(a) = b$. The functor $G : \mathbf{Ring}_* \longrightarrow \mathbf{Ring}$ simply sends (R, a) to R , ‘forgetting’ the specified element a . Show that for each ring R , there is a universal arrow from R to G . (Hint: polynomials!).
2. Assume we have a category \mathcal{C} with a terminal object and binary products. Show that exponentials can be axiomatized in a purely equational fashion, as follows. For each pair of objects A, B , there is an object $A \Rightarrow B$ and a morphism

$$\mathbf{Ap}_{A,B} : (A \Rightarrow B) \times A \longrightarrow B$$

and for each morphism $f : C \times A \longrightarrow B$ there is a morphism $\Lambda(f) : C \longrightarrow (A \Rightarrow B)$ such that

$$f = \mathbf{Ap}_{A,B} \circ (\Lambda(f) \times \text{id}_A).$$

Moreover, for each morphism $g : C \longrightarrow (A \Rightarrow B)$:

$$\Lambda(\mathbf{Ap}_{A,B} \circ (g \times \text{id}_A)) = g.$$

Show that this is equivalent to the definition in terms of couniversal arrows given in the Notes.

3. Let X be a non-empty set. Consider the set H of all those sets of subsets Θ of X ($\Theta \subseteq P(P(X))$) with the following property:

$$U \in \Theta \wedge T \subseteq U \implies T \in \Theta.$$

This set H , ordered by inclusion, is a poset, and hence a category. Show that H is not closed under complements (and in fact is not a Boolean algebra), but that H is Cartesian closed.